

# Large-Scale Knowledge Processing Optimization Techniques (2)



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# Prev. class + $\alpha$ : Linear Programming Problem

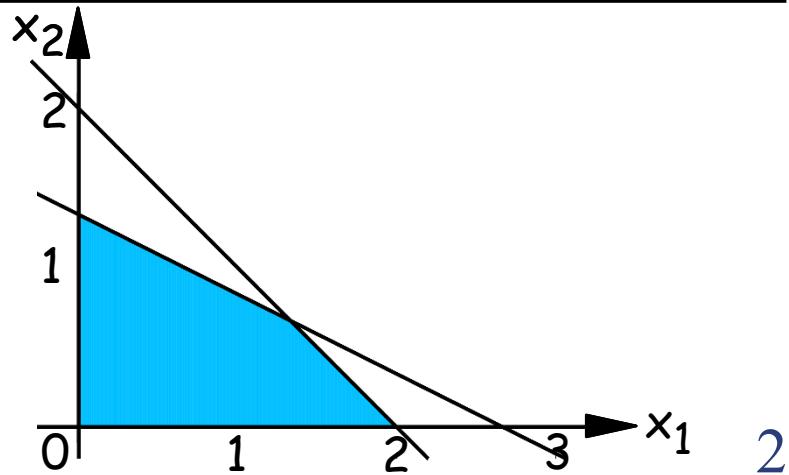
## Standard form

- minimize  $z = -4x_1 - 5x_2$   
subject to  $2x_1 + 2x_2 + x_3 = 4$   
 $3x_1 + 6x_2 + x_4 = 8$   
 $x_1, x_2, x_3, x_4 \geq 0$  Transformation to  
its standard form

Intuitive observation :

- By checking the **extreme points** (vertices) of the feasible region, we can find the optimal solution
- How can we find **extreme points** of the feasible region ?  
(The standard form gives a good hint)

$$\begin{aligned} & \text{maximize } z = 4x_1 + 5x_2 \\ & \text{subject to } 2x_1 + 2x_2 \leq 4 \\ & \quad 3x_1 + 6x_2 \leq 8 \\ & \quad x_1, x_2 \geq 0 \end{aligned}$$



## Ex.) Standard form (Vector representation)

■ minimize  $z = -4x_1 - 5x_2$   
subject to  $2x_1 + 2x_2 + x_3 = 4$   
 $3x_1 + 6x_2 + 4x_4 = 8$   
 $x_j \geq 0 \ (j = 1, 2, \dots, 4)$

■ minimize  $z = (-4 \ -5 \ 0 \ 0) \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix}$

subject to  $\begin{pmatrix} 2 & 2 & 1 & 0 \\ 3 & 6 & 0 & 4 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 4 \\ 8 \end{pmatrix} \geq \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$

## Standard form (Vector representation)

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & & a_{mn} \end{bmatrix}$$

■ minimize  $z = \sum_{j=1}^n c_j x_j$

subject to  $\sum_{j=1}^n a_{ij} x_j = b_i \quad (i = 1, 2, \dots, m)$

$$x_j \geq 0 \quad (j = 1, 2, \dots, n)$$

■ minimize  $z = c^T x$

$$c^T = (c_1, c_2, \dots, c_n)$$

subject to  $A x = b$

$$A = (a_{ij}) \quad m \times n \text{ matrix}$$

$$x \geq 0$$

$$b = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix} \quad x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}$$

## Standard form (Vector representation)

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & & a_{mn} \end{bmatrix}$$

$$\begin{bmatrix} a_{11} \\ a_{21} \\ \vdots \\ a_{m1} \end{bmatrix} \begin{bmatrix} a_{12} \\ a_{22} \\ \vdots \\ a_{m2} \end{bmatrix} \dots \begin{bmatrix} a_{1n} \\ a_{2n} \\ \vdots \\ a_{mn} \end{bmatrix}$$

$A$  can be regarded as  $(A_1 \ A_2 \ \dots \ A_n)$   
(i.e., concatenation of  $n$  vertical vectors)

- minimize  $z = c^T x$   $c^T = (c_1, c_2, \dots, c_n)$   
subject to  $A x = b$   $A = (a_{ij})$   $m \times n$  matrix  
 $x \geq 0$

$$b = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix} \quad x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}$$

## Constraints in p. 3

### Ex.) Basis

- Infinitely many assignments to  $(x_1, x_2, x_3, x_4)$  satisfying the 2 conditions

- In case  $(x_3, x_4) = (0, 0)$  ?
- i.e., if we **focus on  $x_1, x_2$**

$\rightarrow (x_1, x_2, x_3, x_4)$  is uniquely obtained

$$\begin{bmatrix} 2 & 2 & 1 & 0 \\ 3 & 6 & 0 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 4 \\ 8 \end{bmatrix}$$

**Basic vectors:**  $\{ \left( \begin{smallmatrix} 2 \\ 3 \end{smallmatrix} \right), \left( \begin{smallmatrix} 2 \\ 6 \end{smallmatrix} \right) \}$

**Basic matrix:**  $\left( \begin{smallmatrix} 2 & 2 \\ 3 & 6 \end{smallmatrix} \right)$

**Basic variables:**  $x_1, x_2$ , **Nonbasic variables:**  $x_3, x_4$

**Basic solution**

for basic vectors  $\{ \left( \begin{smallmatrix} 2 \\ 3 \end{smallmatrix} \right), \left( \begin{smallmatrix} 2 \\ 6 \end{smallmatrix} \right) \}$ :  $\hat{x} = (4/3, 2/3, 0, 0)$

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \left( \begin{smallmatrix} 2 & 2 \\ 3 & 6 \end{smallmatrix} \right)^{-1} \begin{pmatrix} 4 \\ 8 \end{pmatrix}$$

$$= \begin{pmatrix} 4/3 \\ 2/3 \end{pmatrix}$$

Diagram showing arrows pointing from the text "Nonbasic variables" to  $x_3, x_4$  and from the text "Basic solution" to the vector  $\hat{x}$ .

## Constraints in p. 3

### Ex.) Basis

- Try another base vector

$$\begin{bmatrix} 2 & 2 & 1 & 0 \\ 3 & 6 & 0 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 4 \\ 8 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 2 & 1 & 0 \\ 3 & 6 & 0 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 4 \\ 8 \end{bmatrix}$$

$$\begin{pmatrix} x_1 \\ x_3 \end{pmatrix} = \begin{pmatrix} 2 & 1 \\ 3 & 0 \end{pmatrix}^{-1} \begin{pmatrix} 4 \\ 8 \end{pmatrix}$$

$$= \begin{pmatrix} 8/3 \\ -4/3 \end{pmatrix}$$

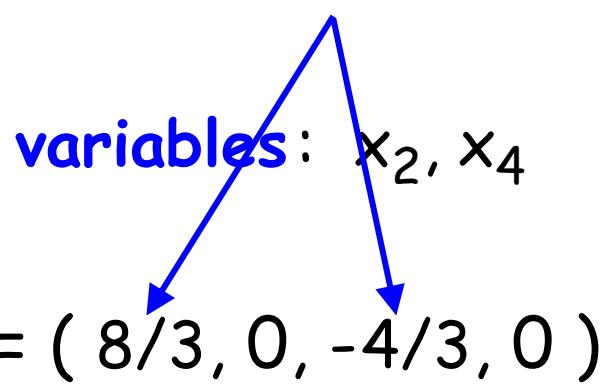
■ **Basic vectors:**  $\{ \begin{pmatrix} 2 \\ 3 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \}$

■ **Basic matrix:**  $\begin{pmatrix} 2 & 1 \\ 3 & 0 \end{pmatrix}$

■ **Basic variables:**  $x_1, x_3$ , **Nonbasic variables:**  $x_2, x_4$

■ **Basic solution**

for basic vectors  $\{ \begin{pmatrix} 2 \\ 3 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \}$ :  $\hat{x} = (8/3, 0, -4/3, 0)$



# Supplementary information

- Some basic matrix **may not have the inverse matrix**

$\begin{bmatrix} 1 & 2 \\ 3 & 6 \end{bmatrix}$  and  $\begin{bmatrix} 1 \\ 3 \end{bmatrix}$  and  $\begin{bmatrix} 2 \\ 6 \end{bmatrix}$  are linear dependent (not independent)

- $A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \dots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix}$

- Assumption 1:  $n \geq m$
- Assumption 2: the rank of  $A$  is  $m$
- Assumption 3: optimal solution exists

We can relax these assumptions (later)

# Basic vector

## Basic matrix

### Basic vector

Their combination  
is not unique

$$\begin{bmatrix} a_{11} \\ a_{21} \\ \vdots \\ a_{m1} \end{bmatrix} \begin{bmatrix} a_{12} \\ a_{22} \\ \vdots \\ a_{m2} \end{bmatrix} \cdots \begin{bmatrix} a_{1n} \\ a_{2n} \\ \vdots \\ a_{mn} \end{bmatrix}$$
$$A = (A_1 \ A_2 \ \cdots \ A_n)$$

- $m$  linear independent vectors in  $A$

$$\text{Basis } \mathcal{B} = \{ A_{i1}, A_{i2}, \dots, A_{im} \}$$

$$\begin{pmatrix} 2 & 2 & 1 & 0 \\ 3 & 6 & 0 & 4 \end{pmatrix}$$



$$\mathcal{B} = \left\{ \begin{pmatrix} 2 \\ 3 \end{pmatrix}, \begin{pmatrix} 2 \\ 6 \end{pmatrix} \right\}$$

### Non basic vector

- Other  $(n - m)$  vectors

### Basic matrix B

- $m \times m$  square matrix with  $m$  basic vectors

$$B = (A_{i1} \ A_{i2} \ \cdots \ A_{im})$$

$$\begin{pmatrix} 2 & 2 \\ 3 & 6 \end{pmatrix}$$

### Nonbasic matrix N

- $m \times (n - m)$  matrix with  $(n - m)$  Nonbasic vectors

# Basic variable

$$x_B = (x_1, x_2)$$

# Basic solution

$$x_N = (x_3, x_4)$$

$$\hat{x} = (8/3, 2/3, 0, 0)$$

## Basic variable

- **m** basic variables  $x_{i1}, x_{i2}, \dots, x_{im}$  corresponding to  $m$  basic vectors  $A_{i1}, A_{i2}, \dots, A_{im}$
- $m$ -dimensional vector  $x_B = (x_{i1}, x_{i2}, \dots, x_{im})$

## Nonbasic variable

- Other  **$(n-m)$**  variables
- $(n-m)$ -dimensional vector  $x_N$

## Basic solution $\hat{x}$ for basis $\mathcal{B} = \{A_{i1}, A_{i2}, \dots, A_{im}\}$

- $x_j = \begin{cases} 0 & (A_j \notin \mathcal{B}) \\ \text{l-th element of } B^{-1} b & (A_j \in \mathcal{B}, x_j = x_{il}) \end{cases}$

Solution of  $Bx = b$

## practice: Basic solution

- minimize  $z = -4x_1 - 5x_2$   
 subject to  $2x_1 + 2x_2 + x_3 = 4$   
 $3x_1 + 6x_2 + x_4 = 8$   
 $x_j \geq 0 \ (j = 1, 2, \dots, 4)$
- $A_1 = \begin{pmatrix} 2 \\ 3 \end{pmatrix}$ ,  $A_2 = \begin{pmatrix} 2 \\ 6 \end{pmatrix}$ ,  $A_3 = \begin{pmatrix} \phantom{2} \\ \phantom{2} \end{pmatrix}$ ,  $A_4 = \begin{pmatrix} \phantom{2} \\ \phantom{2} \end{pmatrix}$
- $A_1, A_2$  independent:  $\text{rank}(A_1 A_2) = m$  (determinant  $\neq 0$ )
- Solve  $(A_1 \ A_2) \begin{pmatrix} 4 \\ 8 \end{pmatrix} = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$ : we have  $x_1 = \phantom{0}$ ,  $x_2 = \phantom{0}$
- Basic solution for  $A_1, A_2$  is  $x^T = \begin{pmatrix} \phantom{0} \\ \phantom{0} \\ \phantom{0} \end{pmatrix}$
- All basic elements (elements of basic solution) are non-negative  $\rightarrow$  **basic feasible solution**

## practice: Basic solution

- From the conditions of the problem in p. 3, we have 4 vectors  $A_1 \sim A_4$ . We have 6 ways for choosing 2 basic vectors  $A_i$  and  $A_j$ . For each of them, find basic solutions.
- From each of the above basic solutions, obtain the coordinate  $(x_1, x_2)$ . Find such coordinates in the graph of the original optimization problem.

Original  
optimization prob.

$$\begin{array}{ll} \text{maximize} & 4x_1 + 5x_2 \\ \text{subject to} & 2x_1 + 2x_2 \leq 4 \\ & 3x_1 + 6x_2 \leq 8 \\ & x_1 \geq 0, x_2 \geq 0 \end{array}$$

# Standard form, basic solution

$$A x = b$$

$$x \geq 0$$

$$\begin{matrix} A_1 & A_2 & A_3 & A_4 \\ \left(\begin{matrix} 2 \\ 3 \end{matrix}\right) & \left(\begin{matrix} 2 \\ 6 \end{matrix}\right) & \left(\begin{matrix} 1 \\ 0 \end{matrix}\right) & \left(\begin{matrix} 0 \\ 1 \end{matrix}\right) \end{matrix}$$

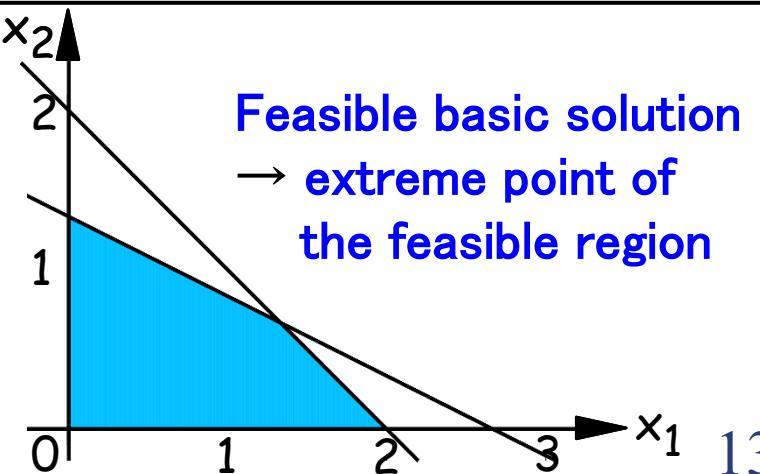
basic solutions for

- $A_1, A_2$
- $A_1, A_3$
- $A_1, A_4$
- $A_2, A_3$
- $A_2, A_4$
- $A_3, A_4$

$\left(\frac{4}{3}, \frac{2}{3}\right)$	$(0, 0)$
$\left(\frac{8}{3}, 0\right)$	$(-\frac{4}{3}, 0)$
$(2, 0)$	$(0, 2)$
$(0, \frac{4}{3})$	$(\frac{4}{3}, 0)$
$(0, 2)$	$(0, -4)$
$(0, 0)$	$(4, 8)$

■ minimize  $z = -4x_1 - 5x_2$   
 subject to  $2x_1 + 2x_2 + x_3 = 4$   
 $3x_1 + 6x_2 + x_4 = 8$   
 $x_1, x_2, x_3, x_4 \geq 0$  Transformation to  
 its standard form

maximize  $z = 4x_1 + 5x_2$   
 subject to  $2x_1 + 2x_2 \leq 4$   
 $3x_1 + 6x_2 \leq 8$   
 $x_1, x_2 \geq 0$



# Degeneration of feasible basic solutions

- Feasible basic solution  $x$  is **degenerate**
  - $x$  has more 0's than  $n-m$
- Two different basis correspond to the same basic solution  $x$ 
  - $x$  is degenerate

## practice: Degeneration

- For the problems in "Optimization Techniques (1)" slides p. 20 (a), p. 21
  - (1) Transform them into their standard forms
  - (2) Check whether they have degenerate feasible basic solution

# Degeneration of feasible basic solutions

■ minimize  $z = -4x_1 - 5x_2$   
subject to  $2x_1 + 2x_2 + x_3 = 4$   
 $3x_1 + 6x_2 + x_4 = 8$   
 $x_1 + 4x_2 + x_5 = 4$   
 $x_1, x_2, x_3, x_4, x_5 \geq 0$

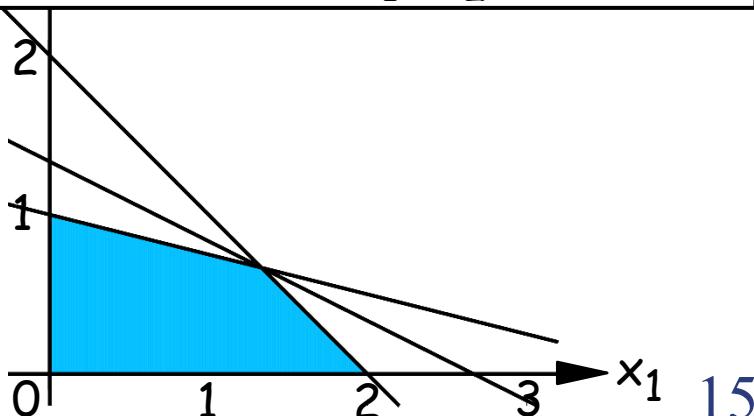
$A_1 \quad A_2 \quad A_3 \quad A_4 \quad A_5$   
 $(\frac{2}{3}) \quad (\frac{2}{6}) \quad (\frac{1}{0}) \quad (\frac{0}{1}) \quad (\frac{0}{1})$

basic solutions for

- $A_1, A_2, A_3 \quad (\frac{4}{3}, \frac{2}{3}, 0, 0, 0)$
- $A_1, A_2, A_4 \quad (\frac{4}{3}, \frac{2}{3}, 0, 0, 0)$
- $A_1, A_2, A_5 \quad (\frac{4}{3}, \frac{2}{3}, 0, 0, 0)$

⋮

maximize  $z = 4x_1 + 5x_2$   
subject to  $2x_1 + 2x_2 \leq 4$   
 $3x_1 + 6x_2 \leq 8$   
 $x_1 + 4x_2 \leq 4$   
 $x_1, x_2 \geq 0$

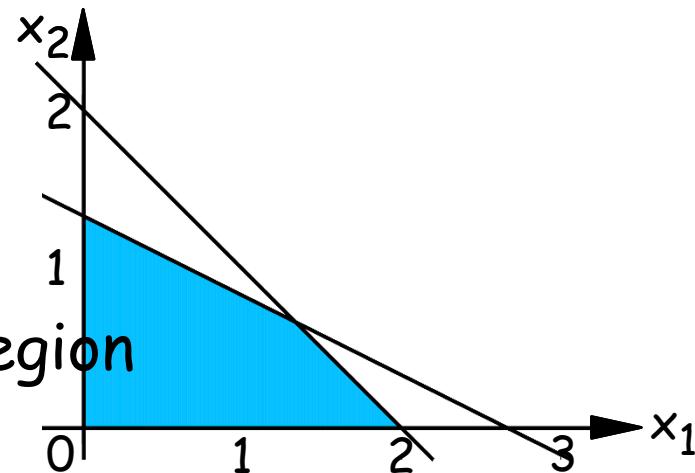


## practice: Basic solution

- Enumerate all basic solutions for the problem in p. 15
  - We have 5 basic vectors  $A_1, A_2, A_3, A_4, A_5$
  - Choose any 3 from these 5

### Summary (1<sup>st</sup> half)

- Feasible basic solution  
→ extreme point of feasible region
- Basis → basic solution
- We can find the optimal solution by checking **all** basic solutions for **all** basis



# Linear programming

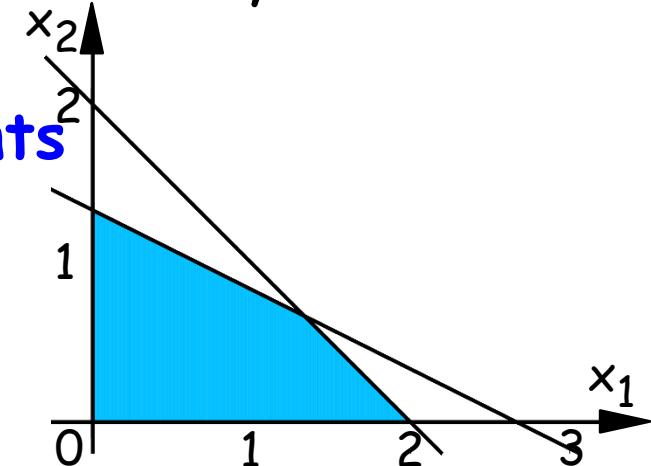
## Simplex method (Simplex algorithm)



- Check all basic solutions ?
- Better ways ?  
Walk through feasible basic solutions  
with improving the objective value

## preliminaries: Feasible region, extreme point

- Lemma: Feasible region  $F$  is convex
- Definition:  $x \in F$  is an **extreme point**
  - The point that cannot be a convex combination of two different points  $x', x'' \in F$
- Properties:
  - The number of **extreme point** for any  $F$  is **finite**
  - $F$  can be represented as a **convex hull of extreme points**



- $S$  is convex
  - $\forall x', x'' \in S, \forall \alpha (0 \leq \alpha \leq 1)$   
 $x' + (1 - \alpha)x'' \in S$
- Convex hull of point set  $S'$ 
  - Minimum convex region containing  $S'$

# Simplex dictionary (Dictionary)

- minimize  $z = -2x_1 + x_2 + x_3 - x_4$   
subject to  $x_1 + 2x_2 - x_3 + x_4 = 0$   
 $2x_1 - 2x_2 + 3x_3 + 4x_4 = 9$   
 $x_1, x_2, x_3, x_4 \geq 0$
- Basic solution for  $A_2 A_3$   
 $(0, \frac{9}{4}, \frac{9}{2}, 0)$  ... feasible solution

## ■ Dictionary

- Solve the conditions for **basic variables**

Current objective value

■  $z = 27/4 - 19/4 x_1 - 21/4 x_4$

■  $x_2 = 9/4 - 5/4 x_1 - 7/4 x_4$

■  $x_3 = 9/2 - 3/2 x_1 - 5/2 x_4$

+ substitute them in the **objective function**

Basic variables      Values in basic solution      Nonbasic variables

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# Basis exchange (Pivot operation)

- $z = 27/4 - 19/4 x_1 - 21/4 x_4$
- $x_2 = 9/4 - 5/4 x_1 - 7/4 x_4$
- $x_3 = 9/2 - 3/2 x_1 - 5/2 x_4$
- $x_1 = x_4 = 0$  holds in the current basic solution
- The coefficients of  $x_1, x_4$  in the objective function are  $-19/4, -21/4$ , respectively
- If we increase  $x_4$  by  $\Delta$ 
  - the objective value decreases:  $27/4 - 21/4 \Delta$
  - If  $x_4$  increases by itself (other variables are fixed), the constraints are not satisfied
    - **change basic variables** ( $x_2, x_3$ )  
**with fixing other nonbasic variables** ( $x_1 = 0$ )

# Basis exchange (Pivot operation)

- $z = 27/4$       -  $21/4 x_4$
- $x_2 = 9/4$       -  $7/4 x_4$
- $x_3 = 9/2$       -  $5/2 x_4$
- **Fix other nonbasic variables ( $x_1 = 0$ )**
- Increase  $x_4$  from 0
  - $x_2$  decreases (when  $x_4 = 9/7$ , we have  $x_2 = 0$ )
  - $x_3$  decreases (when  $x_4 = 9/5$ , we have  $x_3 = 0$ )
- Stop at  $x_4 = 9/7$  to keep the **feasibility** ( $x_i \geq 0$ )
  - $x_2 = 0$ , i.e., " $x_2$  becomes a nonbasic variable"
  - $x_4$  becomes a basic variable

# Basis exchange (Pivot operation)

- $z = 27/4 - 19/4 x_1 - 21/4 x_4$
- $x_2 = 9/4 - 5/4 x_1 - 7/4 x_4$
- $x_3 = 9/2 - 3/2 x_1 - 5/2 x_4$

## Update a dictionary



- $x_2$  becomes a nonbasic variable,  
 $x_4$  becomes a basic variable
- Transform eq.  $x_2 = \dots$  as eq.  $x_4 = \dots$   
& Substitute  $x_4$  in other equations

- $z = -9/4 + 19/5 x_1 + 7/5 x_2$
- $x_4 = 9/7 - 5/7 x_1 - 4/7 x_2$
- $x_3 = 9/5 + 6/5 x_1 - 2/5 x_2$

## practice: Pivot operation

- Update the following simplex dictionary so that  $x_2, x_3$  are basic variables and  $x_1, x_4$  are nonbasic variables

$$z = -4x_1 - 6x_2$$

$$x_3 = 4 - 2x_1 - 2x_2$$

$$x_4 = 9 - 3x_1 - 6x_2$$

# Optimal dictionary

- All coefficients (of nonbasic variables) in the objective function are positive
  - We cannot decrease objective value

Optimal dictionary

Optimal value

- For all feasible solution,  
 $x_2, x_4 \geq 0$  implies  $z \geq -9/4$
- $x^T = (9/5, 0, 9/5, 0)$  is optimal  
with  $z = -9/4$

- $z = -9/4 + 19/5 x_2 + 7/5 x_4$
- $x_1 = 9/5 - 4/5 x_2 - 7/5 x_4$
- $x_3 = 9/5 + 6/5 x_2 - 2/5 x_4$

Optimal solution  $x^T = (9/5, 0, 9/5, 0)$

# Simplex tableau

$$\begin{aligned}
 z &= 27/4 - 19/4 x_1 - 21/4 x_4 \\
 x_2 &= 9/4 - 5/4 x_1 - 7/4 x_4 \\
 x_3 &= 9/2 - 3/2 x_1 - 5/2 x_4
 \end{aligned}$$

## Represent a dictionary by a simplex tableau

Current

objective function

Basic variables

	$x_1$	$x_4$
$z$	27/4	- 19/4 - 21/4
$x_2$	9/4	- 5/4 - 7/4
$x_3$	9/2	- 3/2 - 5/2

Values in basic solution

Nonbasic variables

Relative cost  
coefficients

Rates of change of  
the objective value  
with respect to the  
nonbasic variables

## Pivot operation

- Pivot  $(r, s)$  :  $r$ -th row &  $s$ -th column
- Exchange the  $r$ -th basic var. & the  $s$ -th nonbasic var.
- Ex.: Pivot  $(1, 1)$  to the above tableau

- Transform eq.  $x_2 = \dots$  as eq.  $x_1 = \dots$   
& Substitute  $x_1$  in other equations

Pivot out variable  
(Variable exiting from the basis)

Pivot in variable  
(Variable entering the basis)

# Simplex method

	$x_1$	$x_2$
$z$	0	- 4
$x_3$	4	- 2
$x_4$	9	- 3

$$\begin{aligned} & \text{minimize} \quad z = -4x_1 - 6x_2 \\ & \text{subject to} \quad 2x_1 + 2x_2 + x_3 = 4 \\ & \quad 3x_1 + 6x_2 + x_4 = 9 \\ & \quad x_1, x_2, x_3, x_4 \geq 0 \end{aligned}$$



	$x_1$	$x_4$
$z$	-9	-1 1
$x_3$	1	-1 1/3
$x_2$	3/2	-1/2 -1/6

Pivot (2, 2)

Pivot the 2<sup>nd</sup> row and 2<sup>nd</sup> column



$z$		

Pivot ( , )

# Summary (2<sup>nd</sup> half)

- Represent a feasible basic solution  
(an extreme point of the feasible region)  
by a simplex tableau
- Overview of Simplex method
  - Find a starting point:  
Find a feasible basic solution (given in the next class)
  - Pivot operation:  
Exchange the r-th basic var. & the s-th nonbasic var.  
→ Walk through the extreme points  
(until we cannot improve  
the objective function)

