

Large-Scale Knowledge Processing

# Optimization Techniques (3)



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Walk through the extreme points from an initial feasible basic solution

## Prev. class + $\alpha$ : Simplex method



Initial feasible basic solution

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $z$   | 0     | -4    |
| $x_3$ | 4     | -2    |
| $x_4$ | 8     | -3    |
| $x_5$ | 4     | -1    |

|       | $x_1$ | $x_5$ |
|-------|-------|-------|
| $z$   | -5    | -11/4 |
| $x_3$ | 2     | -3/2  |
| $x_4$ | 2     | -3/2  |
| $x_2$ | 1     | -1/4  |

|       | $x_3$ | $x_5$ |
|-------|-------|-------|
| $z$   | -26/3 | 11/6  |
| $x_1$ | 4/3   | -2/3  |
| $x_4$ | 0     | 1     |
| $x_2$ | 2/3   | 1/6   |

$$\begin{aligned}
 & \text{minimize } z = -4x_1 - 5x_2 \\
 & \text{subject to } 2x_1 + 2x_2 + x_3 = 4 \\
 & \quad 3x_1 + 6x_2 + x_4 = 8 \\
 & \quad x_1 + 4x_2 + x_5 = 4 \\
 & \quad x_1, x_2, x_3, x_4, x_5 \geq 0
 \end{aligned}$$

Pivot  
(3, 2)

$$\begin{aligned}
 z &= -4x_1 - 5x_2 \\
 x_3 &= -2x_1 - 2x_2 + 4 \\
 x_4 &= -3x_1 - 6x_2 + 8 \\
 x_5 &= -x_1 - 4x_2 + 4
 \end{aligned}$$

Increase  $x_2$  from 0  
(Fix other nonbasic var.  $x_1 = 0$ )

Pivot  
(1, 1)

$x_3$  decrease,  $x_3 \geq 0$  if  $x_2 \leq 2$   
 $x_4$  decrease,  $x_4 \geq 0$  if  $x_2 \leq 4/3$   
 $x_5$  decrease,  $x_5 \geq 0$  if  $x_2 \leq 1$

$$x_2 = -1/4x_1 - 1/4x_5 + 1$$

# Initial feasible basic solution ?

- General way will be given later
- Easy case:

$$\begin{array}{ll}\text{maximize} & z = 4x_1 + 6x_2 \\ \text{subject to} & 2x_1 + 2x_2 \leq 4 \\ & 3x_1 + 6x_2 \leq 9 \\ & x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{ll}\text{minimize} & z = -4x_1 - 6x_2 \\ \text{subject to} & 2x_1 + 2x_2 + x_3 = 4 \\ & 3x_1 + 6x_2 + x_4 = 9 \\ & x_1, x_2, x_3, x_4 \geq 0\end{array}$$

Transformation to its standard form

- Basic vars.:  $x_3, x_4$ 
  - $z = -4x_1 - 6x_2$
  - $x_3 = 4 - 2x_1 - 2x_2$
  - $x_4 = 9 - 3x_1 - 6x_2$

## Choice of pivot ?

- One of the typical ways:  
select variable  $x_j$  whose coefficient  $c_j$  is  $c_j < 0$   
and  $|c_j|$  is the largest

# Tips: Pivot operation

just update the simplex tableau

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $z$   | 0     | - 4   |
| $x_3$ | 4     | - 2   |
| $x_4$ | 9     | - 3   |

|       | $x_1$ | $x_4$   |
|-------|-------|---------|
| $z$   |       |         |
| $x_3$ |       |         |
| $x_2$ | $3/2$ | $- 1/2$ |

$9 / -(-6)$        $1 / (-6)$

Ex.: Pivot (2, 2)

Exchange the var. in the 2nd row and the var. in the 2nd column

2nd row

$$x_4 = b_4 + a_1 x_1 + \underline{a_2 x_2}$$

$$x_2 = b_4 / (-a_2) + a_1 / (-a_2) x_1 + 1/a_2 x_4$$

Divide all elements in the 2nd row by  $-a_2$ , but  $1/a_2$  in the 2nd column

# Tips: Pivot operation

just update the simplex tableau

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $z$   | 0     | -4    |
| $x_3$ | 4     | -2    |
| $x_4$ | 9     | -3    |

|       | $x_1$ | $x_4$ |
|-------|-------|-------|
| $z$   |       |       |
| $x_3$ |       |       |
| $x_2$ | (3/2) | -1/2  |

|       | $x_1$ | $x_4$ |
|-------|-------|-------|
| $z$   |       |       |
| $x_3$ | 1     | -1    |
| $x_2$ | 3/2   | -1/2  |

Ex.: Pivot (2, 2)

Exchange the var. in the 2nd row and the var. in the 2nd column

Other rows

$$x_3 = b_3 + a_1 x_1 + \underline{a_2 x_2}$$

$$x_2 = b'_2 + a'_1 x_1 + a'_4 x_4$$



$$(Eq. on x_2) \times a_2 + (Eq. on x_3)$$

i.e., i-th column becomes  $a_i; a_2 + a'_i$

Notice: 2nd column:  $(Eq. x_2) \times a_2$

$$(-1/6)(-2)$$

## Practice: Simplex method

Solve the following linear programming problems by simplex method

- a. "Optimization Techniques (2)" slide p. 24
- b. "Optimization Techniques (2)" slide p. 13
- c. "Optimization Techniques (2)" slide p. 15
- d. "Optimization Techniques (1)" slide p. 7

# Linear programming

## Simplex method



- Can we solve all linear programming problem?

# Situation 1: Feasible region is unbounded

Initial feasible basic solution

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $z$   | 0     | -1    |
| $x_3$ | 1     | 1     |

Original problem

$$\begin{aligned} & \text{minimize } z = -x_1 - 2x_2 \\ & \text{subject to} \\ & \quad x_1 + x_2 \geq -1 \\ & \quad x_1, x_2 \geq 0 \end{aligned}$$

Standard form

$$\begin{aligned} & -x_1 - x_2 + x_3 = 1 \\ & x_1, x_2, x_3 \geq 0 \end{aligned}$$

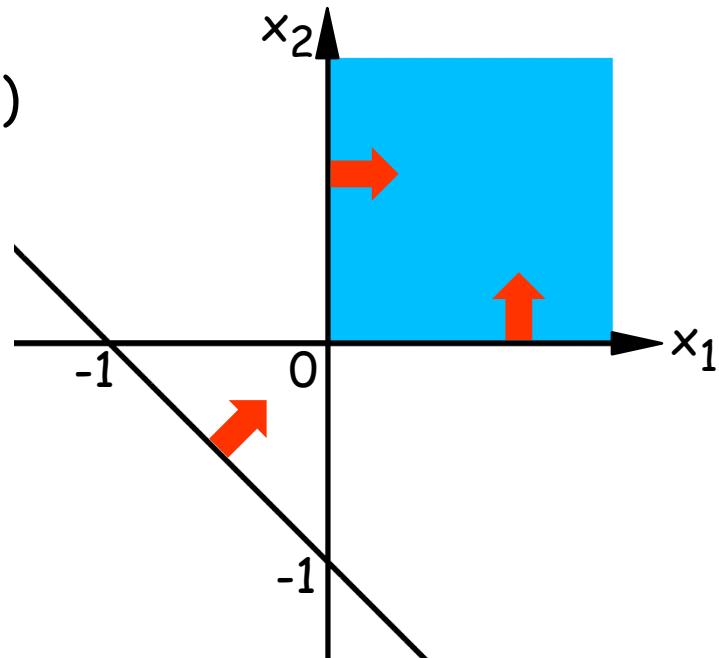
Increase  $x_2$  from 0

(Fix other nonbasic variables as  $x_1 = 0$ )

$x_3$  increases, always satisfies  $x_3 \geq 0$

→ Objective value  $z$  decreases

Feasible region is **unbounded**



## Situation 2: Infeasible (no feasible solutions)

Initial feasible basic solution

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $x_3$ | 0     | -1    |
|       | -1    | -2    |
|       | -1    | -1    |

Original problem

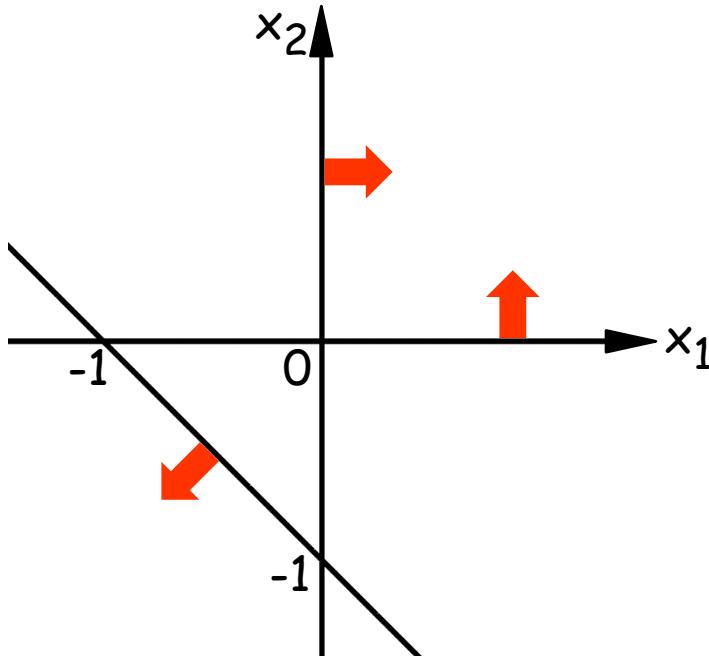
$$\begin{aligned} & \text{minimize } z = -x_1 - 2x_2 \\ & \text{subject to} \quad x_1 + x_2 \leq -1 \\ & \quad x_1, x_2 \geq 0 \end{aligned}$$

Standard form

$$\begin{aligned} & -x_1 - x_2 - x_3 = 1 \\ & x_1, x_2, x_3 \geq 0 \end{aligned}$$

Not a feasible basic solution

How to decide the problem is infeasible ???



## Situation 3: Initial feasible basic solution

Initial feasible basic solution

|       | $x_1$ | $x_2$ |
|-------|-------|-------|
| $z$   | 0     | -1 -2 |
| $x_3$ | -1    | 1 1   |

Original problem

$$\begin{aligned} & \text{minimize } z = -x_1 - 2x_2 \\ & \text{subject to } x_1 + x_2 \geq 1 \\ & \quad x_1, x_2 \geq 0 \end{aligned}$$

Standard form

$$\begin{aligned} & x_1 + x_2 - x_3 = 1 \\ & x_1, x_2, x_3 \geq 0 \end{aligned}$$

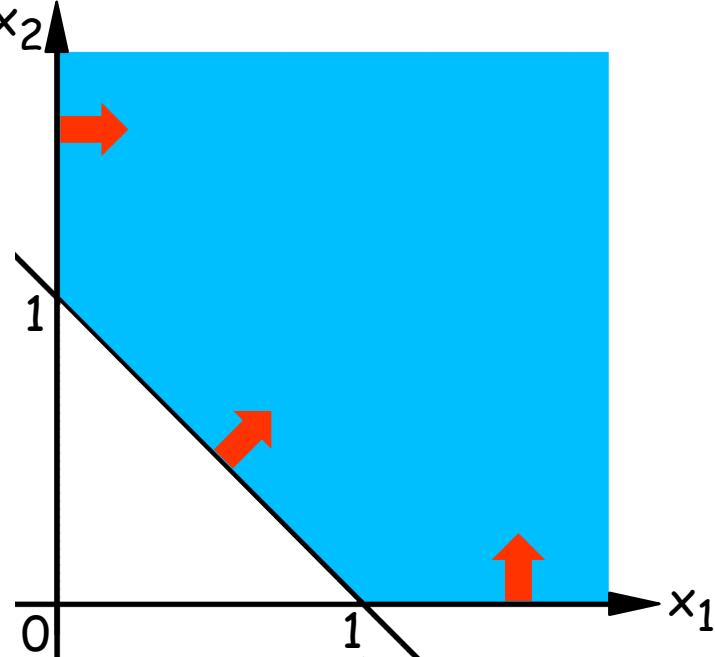
Not a feasible basic solution

(The right fig. says (0, 0) is not a feasible solution, but we have feasible region)

|       | $x_1$ | $x_3$ |
|-------|-------|-------|
| $z$   | 0     | 1 -2  |
| $x_2$ | 1     | -1 1  |

(0, 1) is a feasible basic solution

How to find an initial feasible basic solution ???



# Two-phase simplex method

## Original problem (in standard form)

$$\begin{aligned} \text{minimize } z &= -x_1 - 5x_2 \\ \text{subject to } &4x_1 - x_2 + 4x_3 = 6 \\ &x_1 + 2x_2 + 2x_3 = 4 \\ &x_1, x_2, x_3 \geq 0 \end{aligned}$$

How to find an initial feasible basic solution ???

## Artificial problem

$$\begin{aligned} \text{minimize } w &= x_4 + x_5 \\ \text{subject to } &4x_1 - x_2 + 4x_3 + x_4 = 6 \\ &x_1 + 2x_2 + 2x_3 + x_5 = 4 \\ &x_1, x_2, x_3, x_4, x_5 \geq 0 \end{aligned}$$

Introduce an **artificial variable** for each eq.

Minimize the sum of the **artificial vars.** (so that **all artificial vars.** are **0s.**)

Initial feasible basic solution

$$\begin{aligned} (x_1, x_2, x_3, x_4, x_5) \\ = (0, 0, 0, 6, 4) \end{aligned}$$

Optimal solution of the artificial problem

Phase I: apply simplex method to the artificial problem

$$(0, 2/5, 8/5, 0, 0)$$

This solution satisfies the constraints of the original prob.

= This is the initial feasible basic solution for the original prob.

Phase II: apply simplex method to the original prob.

# Formulation of linear programming problems (additional techniques)



# max{ } function

- Minimize the maximum of two or more objective functions
  - $\text{minimize } \max\{ f(x), g(x), h(x) \}$

→ Introduce a new variable  $z$

$\text{minimize } z$

subject to  $f(x) \leq z, g(x) \leq z, h(x) \leq z$

## Practice:

- $\text{minimize } x_1 + |x_2|$

# max{ } function

- Minimize the maximum of two or more objective functions
  - $\text{minimize } \max\{ f(x), g(x), h(x) \}$

→ Introduce a new variable  $z$

$\text{minimize } z$

subject to  $f(x) \leq z, g(x) \leq z, h(x) \leq z$

## Practice:

- $\text{minimize } x_1 + |x_2|$

→  $\text{minimize } x_1 + x_2'$   
subject to  $x_2 \leq x_2', -x_2 \leq x_2'$

$$|x_2| = \max\{ x_2, -x_2 \}$$

## Practice: Constraints

(Mixed integer  
programming prob.)

- We have 2 variables  $x$  and  $y$ , where  $x = 0, 1$ , and  $y$  is a real number with  $y \geq 0$ . Show constraint inequalities for each of the following constraints

1.  $y \geq 30$  when  $x = 1$
2.  $y = 0$  when  $x = 0$
3.  $y \leq 10$  when  $x = 0$
4.  $y \leq 10$  when  $x = 0$ , and  $y \geq 30$  when  $x = 1$
5.  $y \leq 10$  when  $x = 0$ , and  $y \geq 30$  when  $x = 1$

# Constraints (Mixed integer programming prob.)

1.  $y \geq 30$  when  $x = 1$

- $y \geq 30x$

i.e.,  $-30x + y \geq 0$

- When  $x = 0$ , this inequality says  $y \geq 0$   
(thus, we have no constraint in this case)
- When  $x = 1$ , we add constraint  $y \geq 30$

# Constraints (Mixed integer programming prob.)

2.  $y = 0$  when  $x = 0$

- $y \leq Mx$  ( $M$  is a sufficiently large constant)  
i.e.,  $-Mx + y \leq 0$
- When  $x = 0$ , we have  $y \leq 0$ .  
As we have  $y \geq 0$ , we have constraint  $y = 0$

3.  $y \leq 10$  when  $x = 0$

- $-Mx + y \leq 10$  ( $M$  is a sufficiently large constant)

# Constraints (Mixed integer programming prob.)

4.  $y \leq 10$  when  $x = 0$ ,  $y \geq 30$  when  $x = 1$ 
  - $-Mx + y \leq 10$ , ( $M$  is a sufficiently large constant)
  - $-30x + y \geq 0$
  - From the above constraints 1 and 3
5.  $y \leq 10$  when  $x = 0$ ,  $y \leq 30$  when  $x = 1$ 
  - $y \leq 10(1 - x) + 30x$   
i.e.,  $-20x + y \leq 10$

## Practice: Formulation

We are given  $n$  points  $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$  in the plane. Each point is colored in red or blue.

1. We'd like to obtain a parabola that is convex at the bottom, such that **all red points** are **above** the parabola and **all blue points** are **below** the parabola. Formulate the problem of **finding the equation of such a parabola** (if it exists), as a linear programming problem.
2. We'd like to obtain a circle such that **all red points** are **inside the circle** (including the boundary) and **all blue points** are **outside the circle** (including the boundary) Formulate the problem of **determining** whether such a **circle exists** or not, as a linear programming problem.

# Formulation (1)

- Without loss of generality, we have  $k$  red points  $(x_1, y_1), \dots, (x_k, y_k)$ , and  $n - k$  blue points  $(x_{k+1}, y_{k+1}), \dots, (x_n, y_n)$
- Let  $y = a x^2 + b x + c$  denote the parabola
- Formulation is as follows:
  - minimize  $0$
  - subject to

Determine whether we have a feasible solution  $(a, b, c)$  satisfying the constraints

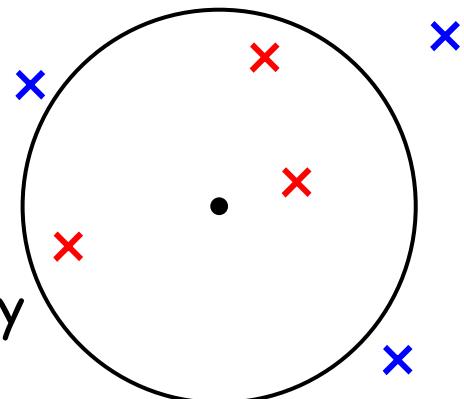
Red points:  $x_i^2 a + x_i b + c \geq y_i \quad (i = 1, 2, \dots, k)$

Blue points:  $x_i^2 a + x_i b + c \leq y_i \quad (i = k+1, k+2, \dots, n)$

$a > 0, b, c$ : free variables

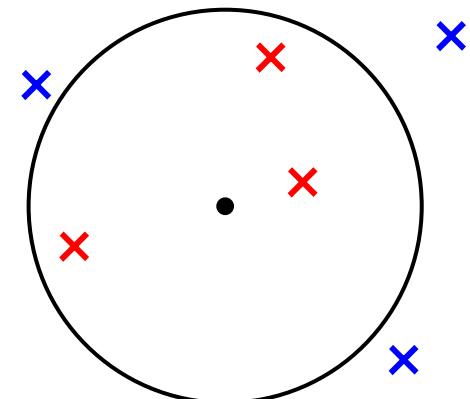
## Formulation (2)

- Without loss of generality, we have  $k$  red points  $(x_1, y_1), \dots, (x_k, y_k)$ , and  $n - k$  blue points  $(x_{k+1}, y_{k+1}), \dots, (x_n, y_n)$
- Let  $(x, y)$  and  $r$  denote the center and the radius of the circle, respectively
- Constraints for each point:  
Red points:  $(x_i - x)^2 + (y_i - y)^2 \leq r^2 \quad (i = 1, 2, \dots, k)$   
Blue points:  $(x_i - x)^2 + (y_i - y)^2 \geq r^2 \quad (i = k+1, k+2, \dots, n)$
- As we have  $x^2, y^2, r^2$  in the constraints, we cannot formulate the problem as a linear programming problem ... (is that right ?)



## Formulation (2)

- Let point  $C(x, y)$  be the center of the circle
- We try to denote the constraints in another way
- Any pair of **red point  $P_i(x_i, y_i)$**  and **blue point  $P_j(x_j, y_j)$**  satisfies **distance  $CP_i \leq$  distance  $CP_j$**
- That is, for any  $i = 1, 2, \dots, k$  and  $j = k+1, k+2, \dots, n$ , we have a constraint
$$(x_i - x)^2 + (y_i - y)^2 \leq (x_j - x)^2 + (y_j - y)^2$$
- From this inequality, we have a constraint
$$2(x_i - x)x + 2(y_i - y)y \geq x_i^2 + y_i^2 - x_j^2 - y_j^2$$



## Formulation (2)

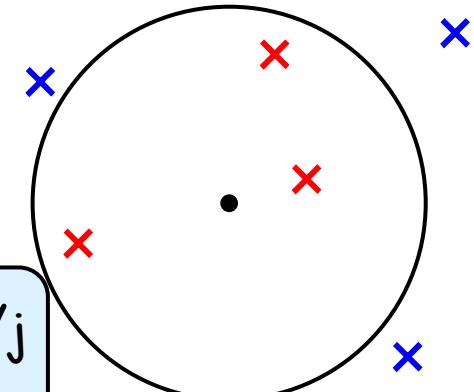
- Formulated linear programming problem is as follows:

Determine whether we have a feasible solution  $(a, b, c)$  satisfying the constraints

- minimize  $0$
- subject to

Note that  $x_i, y_i, x_j, y_j$  are constants

$$2(x_i - x_j)x + 2(y_i - y_j)y \geq x_i^2 + y_i^2 - x_j^2 - y_j^2$$
$$(i = 1, 2, \dots, k; j = k+1, k+2, \dots, n)$$



- By solving this linear programming problem, we can obtain the center  $(x, y)$  of the circle. The radius is not obtained.

# Summary

- Simplex method
  - How to find an initial feasible basic solution ?
  - How to select a pivot ?
- Situations we need to consider
  - Two-phase simplex method
- Formulation of linear programming problems