

Large-Scale Knowledge Processing

# Optimization Techniques (3)

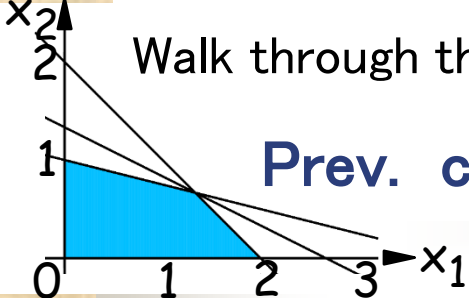


Faculty of Information Science  
and Technology, Hokkaido Univ.

Takashi Horiyama

Walk through the extreme points from an initial feasible basic solution

## Prev. class + $\alpha$ : Simplex method



Initial feasible basic solution

		$x_1$	$x_2$
$z$	0	-4	-5
$x_3$	4	-2	-2
$x_4$	8	-3	-6
$x_5$	4	-1	-4

		$x_1$	$x_5$
$z$	-5	-11/4	5/4
$x_3$	2	-3/2	1/2
$x_4$	2	-3/2	-3/2
$x_2$	1	-1/4	-1/4

		$x_3$	$x_5$
$z$	-26/3	11/6	1/3
$x_1$	4/3	-2/3	1/3
$x_4$	0	1	1
$x_2$	2/3	1/6	-1/3

Pivot  
(3, 2)

Pivot  
(1, 1)

$$\begin{aligned}
 &\text{minimize } z = -4x_1 - 5x_2 \\
 &\text{subject to } \begin{aligned} 2x_1 + 2x_2 + x_3 &= 4 \\ 3x_1 + 6x_2 + x_4 &= 8 \\ x_1 + 4x_2 + x_5 &= 4 \\ x_1, x_2, x_3, x_4, x_5 &\geq 0 \end{aligned}
 \end{aligned}$$

$$\begin{aligned}
 z &= -4x_1 - 5x_2 \\
 x_3 &= -2x_1 - 2x_2 + 4 \\
 x_4 &= -3x_1 - 6x_2 + 8 \\
 x_5 &= -x_1 - 4x_2 + 4
 \end{aligned}$$

Increase  $x_2$  from 0  
(Fix other nonbasic var.  $x_1 = 0$ )

$$\begin{aligned}
 x_3 \text{ decrease, } x_3 &\geq 0 \text{ if } x_2 \leq 2 \\
 x_4 \text{ decrease, } x_4 &\geq 0 \text{ if } x_2 \leq 4/3 \\
 x_5 \text{ decrease, } x_5 &\geq 0 \text{ if } x_2 \leq 1
 \end{aligned}$$

$$x_2 = -1/4 x_1 - 1/4 x_5 + 1$$

# Initial feasible basic solution ?

- General way will be given later
- Easy case:

$$\begin{array}{ll}\text{maximize} & z = 4x_1 + 6x_2 \\ \text{subject to} & 2x_1 + 2x_2 \leq 4 \\ & 3x_1 + 6x_2 \leq 9 \\ & x_1, x_2 \geq 0\end{array}$$



$$\begin{array}{ll}\text{minimize} & z = -4x_1 - 6x_2 \\ \text{subject to} & 2x_1 + 2x_2 + x_3 = 4 \\ & 3x_1 + 6x_2 + x_4 = 9 \\ & x_1, x_2, x_3, x_4 \geq 0\end{array}$$

Transformation to its standard form

- Basic vars.:  $x_3, x_4$ 
  - $z = -4x_1 - 6x_2$
  - $x_3 = 4 - 2x_1 - 2x_2$
  - $x_4 = 9 - 3x_1 - 6x_2$

## Choice of pivot ?

- One of the typical ways:  
select variable  $x_j$  whose coefficient  $c_j$  is  $c_j < 0$   
and  $|c_j|$  is the largest

# Tips: Pivot operation

just update the  
simplex tableau

Ex.: Pivot (**2**, 2)

Exchange the var. in the **2nd row**  
and the var. in the 2nd column

	$x_1$	$x_2$
$z$	0	- 4
$x_3$	4	- 2
$x_4$	9	- 6

**2nd row**

$$x_4 = b_4 + a_1 x_1 + \underline{a_2 x_2}$$

$$x_2 = b_4 / (-a_2) + a_1 / (-a_2) x_1 + 1/a_2 x_4$$

Divide all elements in the **2nd row**  
by  $-a_2$ , but  $1/a_2$  in the 2nd column

	$x_1$	$x_4$
$z$		
$x_3$		
$x_2$	3/2	- 1/6

$$9 / -(-6)$$

$$1 / (-6)$$

# Tips: Pivot operation

just update the  
simplex tableau

	$x_1$	$x_2$
z	0	-6
$x_3$	4	-2
$x_4$	9	-6

Ex.: Pivot (2, 2)

Exchange the var. in the 2nd row  
and the var. in the 2nd column

Other rows

$$x_3 = b_3 + a_1 x_1 + \underline{a_2 x_2}$$

$$x_2 = b'_2 + a'_1 x_1 + a'_4 x_4$$



(Eq. on  $x_2$ )  $\times a_2$  + (Eq. on  $x_3$ )

i.e., i-th column becomes  $a_i a_2 + a'_i$

Notice: 2nd column: (Eq.  $x_2$ )  $\times a_2$

	$x_1$	$x_4$
z		
$x_3$		
$x_2$	3/2	-1/6

$$4 + (3/2)(-2)$$

	$x_1$	$x_4$
z		
$x_3$	1	1/3
$x_2$	3/2	-1/6

$$(-1/6)(-2)$$

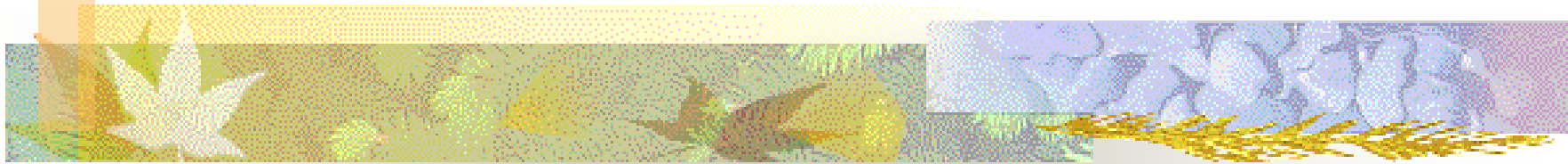
## Practice: Simplex method

Solve the following linear programming problems by simplex method

- a. "Optimization Techniques (2)" slide p. 24
- b. "Optimization Techniques (2)" slide p. 13
- c. "Optimization Techniques (2)" slide p. 15
- d. "Optimization Techniques (1)" slide p. 7

# Linear programming

## Simplex method



- Can we solve all linear programming problem ?

# Situation 1: Feasible region is unbounded

Initial feasible basic solution

		$x_1$	$x_2$
$z$	0	-1	-2
$x_3$	1	1	1

Original problem

$$\begin{aligned} &\text{minimize } z = -x_1 - 2x_2 \\ &\text{subject to } x_1 + x_2 \geq -1 \\ &\quad x_1, x_2 \geq 0 \end{aligned}$$

Standard form

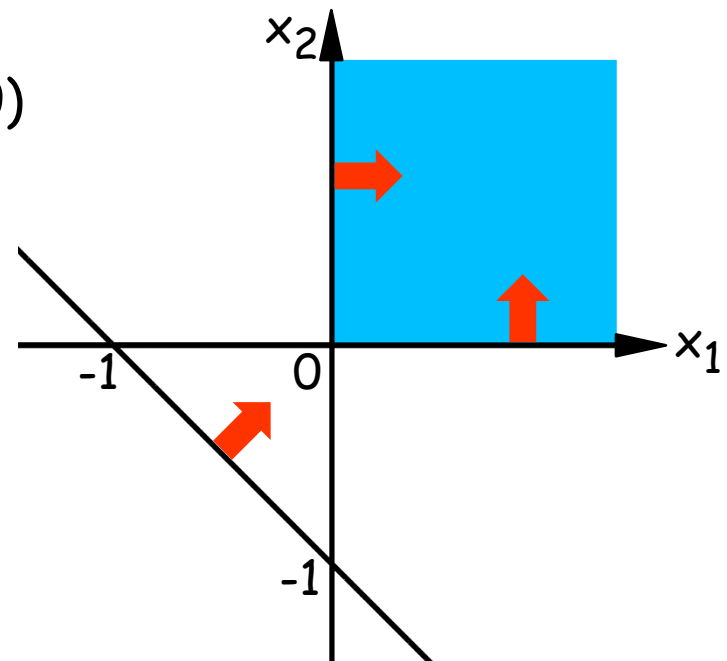
$$\begin{aligned} &-x_1 - x_2 + x_3 = 1 \\ &x_1, x_2, x_3 \geq 0 \end{aligned}$$

Increase  $x_2$  from 0  
(Fix other nonbasic variables as  $x_1 = 0$ )

$x_3$  increases, always satisfies  $x_3 \geq 0$

→ Objective value  $z$  decreases

Feasible region is **unbounded**





## Situation 2: Infeasible (no feasible solutions)

Initial feasible basic solution

	$x_1$	$x_2$
$z$	0	-2
$x_3$	-1	-1

Not a feasible basic solution

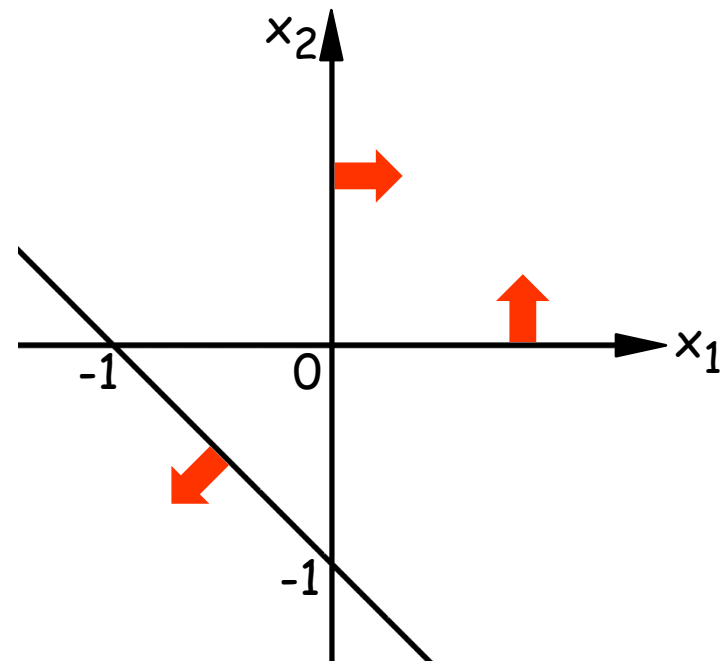
How to decide the problem is infeasible ???

Original problem

$$\begin{aligned} \text{minimize } z &= -x_1 - 2x_2 \\ \text{subject to } x_1 + x_2 &\leq -1 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Standard form

$$\begin{aligned} -x_1 - x_2 - x_3 &= 1 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$



## Situation 3: Initial feasible basic solution

Initial feasible basic solution

	$x_1$	$x_2$
$z$	0	-1 -2
$x_3$	-1	1 1

Not a feasible basic solution

(The right fig. says  $(0, 0)$  is not a feasible solution, but we have feasible region)

	$x_1$	$x_3$
$z$	0	1 -2
$x_2$	1	-1 1

$(0, 1)$  is a feasible basic solution

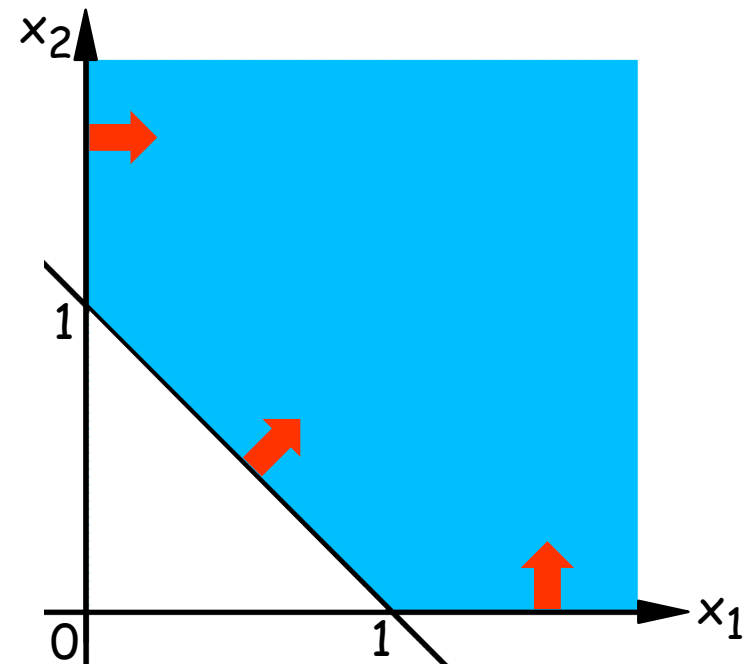
How to find an initial feasible basic solution ???

Original problem

$$\begin{aligned} \text{minimize } z &= -x_1 - 2x_2 \\ \text{subject to } x_1 + x_2 &\geq 1 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Standard form

$$\begin{aligned} x_1 + x_2 - x_3 &= 1 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$



# Two-phase simplex method

## ■ Original problem (in standard form)

$$\begin{aligned} \text{minimize } z &= -x_1 - 5x_2 \\ \text{subject to } 4x_1 - x_2 + 4x_3 &= 6 \\ x_1 + 2x_2 + 2x_3 &= 4 \\ x_1, x_2, x_3 &\geq 0 \end{aligned}$$

How to find an initial feasible basic solution ???

## ■ Artificial problem

$$\begin{aligned} \text{minimize } w &= x_4 + x_5 \\ \text{subject to } 4x_1 - x_2 + 4x_3 + x_4 &= 6 \\ x_1 + 2x_2 + 2x_3 + x_5 &= 4 \\ x_1, x_2, x_3, x_4, x_5 &\geq 0 \end{aligned}$$

Introduce an **artificial variable** for each eq.

Minimize the sum of the **artificial vars.**  
(so that **all artificial vars.** are **0s.**)

Optimal solution of the artificial problem

Phase I: apply simplex method to the artificial problem

$(0, 2/5, 8/5, 0, 0)$

Initial feasible basic solution

$(x_1, x_2, x_3, x_4, x_5)$   
 $= (0, 0, 0, 6, 4)$

This solution satisfies the constraints of the original prob.  
= This is the initial feasible basic solution for the original prob.

Phase II: apply simplex method to the original prob.

# Formulation of linear programming problems (additional techniques)



# max{ } function

- Minimize the maximum of two or more objective functions

- minimize  $\max\{f(x), g(x), h(x)\}$

→ Introduce a new variable  $z$

minimize  $z$

subject to  $f(x) \leq z, g(x) \leq z, h(x) \leq z$

## Practice:

- minimize  $x_1 + |x_2|$

# max{ } function

- Minimize the maximum of two or more objective functions

- minimize  $\max\{f(x), g(x), h(x)\}$


→ Introduce a new variable  $z$

minimize  $z$

subject to  $f(x) \leq z, g(x) \leq z, h(x) \leq z$

## Practice:

- minimize  $x_1 + |x_2|$


$$|x_2| = \max\{x_2, -x_2\}$$

→ minimize  $x_1 + x_2'$   
subject to  $x_2 \leq x_2', -x_2 \leq x_2'$

## Practice: Constraints

(Mixed integer programming prob.)

- We have 2 variables  $x$  and  $y$ , where  $x = 0, 1$ , and  $y$  is a real number with  $y \geq 0$ . Show constraint inequalities for each of the following constraints
- 1.  $y \geq 30$  when  $x = 1$
- 2.  $y = 0$  when  $x = 0$
- 3.  $y \leq 10$  when  $x = 0$
- 4.  $y \leq 10$  when  $x = 0$ , and  $y \geq 30$  when  $x = 1$
- 5.  $y \leq 10$  when  $x = 0$ , and  $y \leq 30$  when  $x = 1$

# Constraints (Mixed integer programming prob.)

1.  $y \geq 30$  when  $x = 1$

- $y \geq 30x$

i.e.,  $-30x + y \geq 0$

- When  $x = 0$ , this inequality says  $y \geq 0$  (thus, we have no constraint in this case)
- When  $x = 1$ , we add constraint  $y \geq 30$



# Constraints (Mixed integer programming prob.)

2.  $y = 0$  when  $x = 0$

- $y \leq Mx$  ( $M$  is a sufficiently large constant)  
i.e.,  $-Mx + y \leq 0$

- When  $x = 0$ , we have  $y \leq 0$ .

As we have  $y \geq 0$ , we have constraint  $y = 0$

3.  $y \leq 10$  when  $x = 0$

- $-Mx + y \leq 10$  ( $M$  is a sufficiently large constant)

# Constraints (Mixed integer programming prob.)

4.  $y \leq 10$  when  $x = 0$ ,  $y \geq 30$  when  $x = 1$
- $-Mx + y \leq 10$ , ( $M$  is a sufficiently large constant)
  - $-30x + y \geq 0$
  - From the above constraints 1 and 3
5.  $y \leq 10$  when  $x = 0$ ,  $y \leq 30$  when  $x = 1$
- $y \leq 10(1 - x) + 30x$   
i.e.,  $-20x + y \leq 10$

## Practice: Formulation

From Naoki Katoh, Mathematical programming, Corona publishing, p. 58 [7]

We are given  $n$  points  $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$  in the plane. Each point is colored in red or blue.

1. We'd like to obtain a parabola that is convex at the bottom, such that **all red points** are **above** the parabola and **all blue points** are **below** the parabola. Formulate the problem of **finding the equation of such a parabola** (if it exists), as a linear programming problem.
2. We'd like to obtain a circle such that all **red points** are **inside the circle** (including the boundary) and all **blue points** are **outside the circle** (including the boundary). Formulate the problem of **determining** whether such a **circle exists** or not, as a linear programming problem.

# Formulation (1)

- Without loss of generality, we have  
k red points  $(x_1, y_1), \dots, (x_k, y_k)$ , and  
n - k blue points  $(x_{k+1}, y_{k+1}), \dots, (x_n, y_n)$
- Let  $y = a x^2 + b x + c$  denote the parabola
- Formulation is as follows:

- minimize 0

- subject to

Determine whether we have a feasible solution  $(a, b, c)$  satisfying the constraints

Red points:  $x_i^2 a + x_i b + c \geq y_i \quad (i = 1, 2, \dots, k)$

Blue points:  $x_i^2 a + x_i b + c \leq y_i \quad (i = k+1, k+2, \dots, n)$

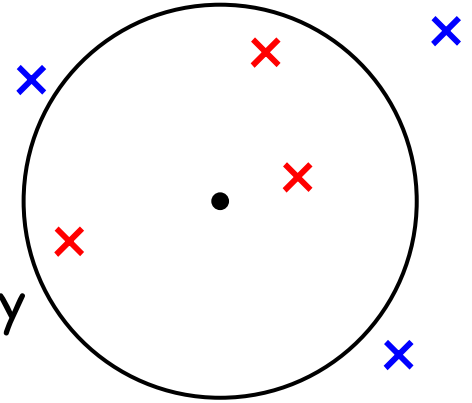
$a > 0$ ,  $b, c$ : free variables

# Formulation (2)

From Naoki Katoh, Mathematical programming, Corona publishing, p. 58 [7]

- Without loss of generality, we have  $k$  red points  $(x_1, y_1), \dots, (x_k, y_k)$ , and  $n - k$  blue points  $(x_{k+1}, y_{k+1}), \dots, (x_n, y_n)$

- Let  $(x, y)$  and  $r$  denote the center and the radius of the circle, respectively



- Constraints for each point:

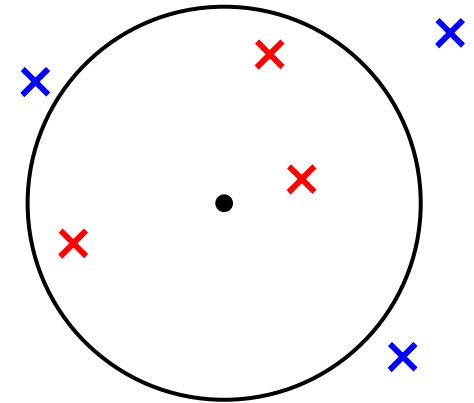
Red points:  $(x_i - x)^2 + (y_i - y)^2 \leq r^2 \quad (i = 1, 2, \dots, k)$

Blue points:  $(x_i - x)^2 + (y_i - y)^2 \geq r^2 \quad (i = k+1, k+2, \dots, n)$

- As we have  $x^2, y^2, r^2$  in the constraints, we cannot formulate the problem as a linear programming problem ... (is that right ?)

## Formulation (2)

- Let point  $C(x, y)$  be the center of the circle
- We try to denote the constraints in another way
- Any pair of **red point**  $P_i(x_i, y_i)$  and **blue point**  $P_j(x_j, y_j)$  satisfies **distance  $C P_i \leq$  distance  $C P_j$**
- That is, for any  $i = 1, 2, \dots, k$  and  $j = k+1, k+2, \dots, n$ , we have a constraint
$$(x_i - x)^2 + (y_i - y)^2 \leq (x_j - x)^2 + (y_j - y)^2$$
- From this inequality, we have a constraint
$$2(x_i - x_j)x + 2(y_i - y_j)y \geq x_i^2 + y_i^2 - x_j^2 - y_j^2$$



# Formulation (2)

From Naoki Katoh, Mathematical programming, Corona publishing, p. 58 [7]

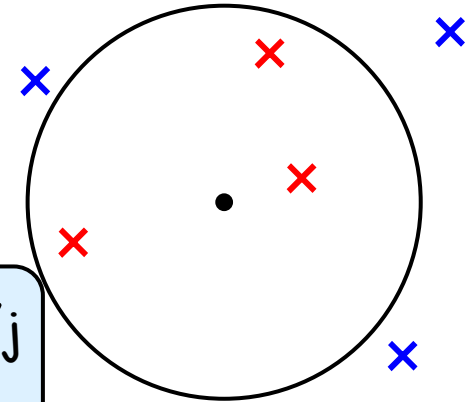
- Formulated linear programming problem is as follows:

Determine whether we have a feasible solution  $(a, b, c)$  satisfying the constraints

- minimize 0
- subject to

Note that  $x_i, y_i, x_j, y_j$  are constants

$$2(x_i - x_j)x + 2(y_i - y_j)y \geq x_i^2 + y_i^2 - x_j^2 - y_j^2 \\ (i = 1, 2, \dots, k; j = k+1, k+2, \dots, n)$$



※ By solving this linear programming problem, we can obtain the center  $(x, y)$  of the circle. The radius is not obtained.



# Summary

- Simplex method
  - How to find an initial feasible basic solution ?
  - How to select a pivot ?
- Situations we need to consider
  - Two-phase simplex method
- Formulation of linear programming problems